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Research Article

A Seed-Based Plant Propagation Algorithm: The Feeding Station Model

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The seasonal production of fruit and seeds is akin to opening a feeding station, such as a restaurant. Agents coming to feed on the fruit are like customers attending the restaurant; they arrive at a certain rate and get served at a certain rate following some appropriate processes. The same applies to birds and animals visiting and feeding on ripe fruit produced by plants such as the strawberry plant. This phenomenon underpins the seed dispersion of the plants. Modelling it as a queuing process results in a seed-based search/optimisation algorithm. This variant of the Plant Propagation Algorithm is described, analysed, tested on nontrivial problems, and compared with well established algorithms. The results are included.

1. Introduction

Plants have evolved a variety of ways to propagate. Propagation with seeds is perhaps the most common of them all and one which takes advantage of all sorts of agents ranging from wind to water, birds, and animals. In [1] a Plant Propagation Algorithm based on the way the strawberry plant propagates using runners has been introduced. Here, we consider the case where the strawberry plant uses seeds to propagate.

Plants rely heavily on the dispersion of their seeds to colonise new territories and to improve their survival [2, 3]. There are a lot of studies and models of seed dispersion particularly for trees [2–6]. Dispersion by wind and ballistic means is probably the most studied of all approaches [7–9]. However, in the case of the strawberry plant, given the way the seeds stick to the surface of the fruit (Figure 1(a)) [10], dispersion by wind or mechanical means is very limited. Animals, however, and birds in particular are the ideal agents for dispersion [2, 3, 11, 12] in this case.

There are many biologically inspired optimization algorithms in the literature [13, 14]. The Flower Pollination Algorithm (FPA) is inspired by the pollination of flowers through different agents [8]; the swarm data clustering algorithm is inspired by pollination by bees [15]; Particle Swarm Optimization (PSO) is inspired by the foraging behavior of

groups of animals and insects [16, 17]; the Artificial Bee Colony (ABC) simulates the foraging behavior of honey bees [18, 19]; the Firefly algorithm is inspired by the flashing fireflies when trying to attract a mate [20, 21]; the Social Spider Optimization (SSO) algorithm is inspired by the cooperative behavior of social spiders [22]. The list could easily be extended.

The Plant Propagation Algorithm (PPA) also known as the strawberry algorithm was inspired by the way plants and specifically the strawberry plants propagate using runners [1, 23]. The attraction of PPA is that it can be implemented easily for all sorts of optimization problems. Moreover, it has few algorithm specific arbitrary parameters. It follows the principle that plants in good spots with plenty of nutrients will send many short runners. They send few long runners when in nutrient poor spots. With long runners PPA tries to explore the search space while short runners enable it to exploit the solution space well. In this paper, we investigate an alternative PPA which is entirely based on the propagation by seeds of the strawberry plant. Because of the periodic nature of fruit and seed production, it amounts to setting up a feeding station for the attention of potential seed-dispersing agents [24], Hence the feeding station model used here and the resulting Seed-Based Plant Propagation Algorithm or SbPPA.

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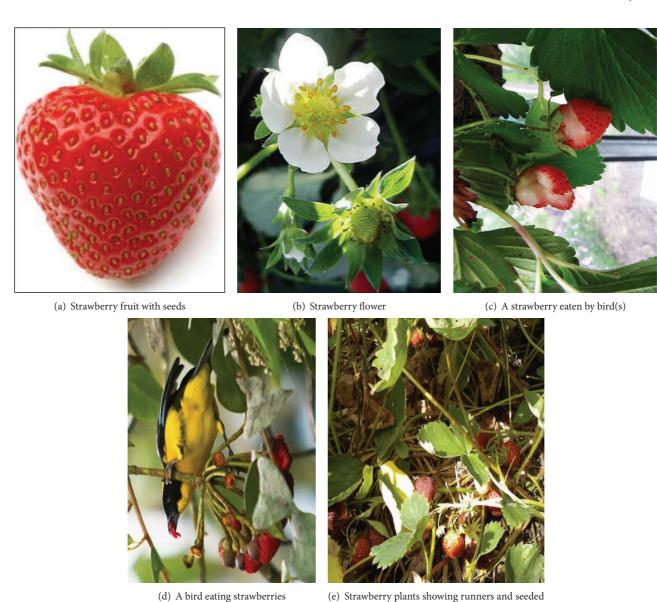


FIGURE 1: Strawberry plant propagation: through seed dispersion [25–28].

SbPPA is tested on both unconstrained and constrained benchmark problems also used in [22, 29, 30]. Experimental results are presented in Tables 4–7 in terms of best, mean, worst, and standard deviation for all algorithms. The paper is organised as follows. In Section 2 we briefly introduce the feeding station model representing strawberry plants in fruit and the main characteristics of the paths followed by different agents that disperse the seeds. Section 3 presents the SbPPA in pseudocode form. The experimental settings, results, and convergence graphs for different problems are given in Section 4.

2. Aspects of the Feeding Station Model

Some animals and plants depend on each other to conserve their species [31]. Thus, many plants require, for effective seed dispersal, the visits of frugivorous birds or animals according to a certain distribution [2, 3, 32, 33].

Seed dispersal by different agents is also called "seed shadow" [32]; this shows the abundance of seeds spread globally or locally around parent plants. Here a queuing model is used which, in the context of a strawberry feeding station model, involves two parts:

- (1) the quantity of fruit or seeds available to agents which implies the rate at which the agents will visit the plants,
- (2) a probability density function that tells us about the service rate with which the agents are served by the plants.

The model estimates the quantity of seeds that is spread locally compared to that dispersed globally [34–38]. There are two aspects that need to be balanced: exploitation, which

is represented by the dispersal of seeds around the plants, and exploration which ensures that the search space is well covered.

Agents arrive at plants in a random process. Assume that at most one agent arrives to the plants in any unit of time (orderliness condition). It is further supposed that the probability of arrivals of agents to the plants remains the same for a particular period of time. This period corresponds to when the plants are in fruit and during which time the number of visitors is stable (stationarity condition). Furthermore, it is assumed that the arrival of one agent does not affect the rest of arrivals (independence).

With these assumptions in mind, the arrival of agents to plants follows a Poisson process [39, 40], which can be formally described as follows. Let X' be the random variable representing the number of arrivals per unit of time t. Then, the probability of k arrivals over t is

$$P(X' = k) = \frac{(\lambda t)^k e^{-\lambda t}}{k!},\tag{1}$$

where λ denotes the mean arrival rate of agents per time unit t. On the other hand, the time taken by agents in successfully eating fruit and leaving to disperse its seeds, in other words the service time for agents, is expressed by a random variable which follows the exponential probability distribution [41]. This can be expressed as follows:

$$S(t) = \mu e^{-\mu t}, \tag{2}$$

where μ is the average number of agents that can feed at time t. Let us assume that the arrival rate of agents is less than the fruits available on all plants per unit of time; therefore $\lambda < \mu$.

We assume that the system is in steady state. Let A denote the average number of agents in the strawberry field (some already eating and the rest waiting to feed) and A_q the average number of agents waiting to get the chance to feed. If we denote the average number of agents eating fruits by λ/μ , then by Little's formula [42], we have

$$A = A_q + \frac{\lambda}{\mu}. (3)$$

Since the plant needs to maximise dispersion, this is equivalent to having a large A_q in (3). Therefore, from this equation, we need to solve the following problem:

Maximize
$$A_q = A - \frac{\lambda}{\mu}$$
,
subject to $g_1(\lambda, \mu) = \lambda < \mu + 1$, $\lambda > 0$, $\mu > 0$,

where A=10, which represents the population size in the implementation. The simple limits on the variables are $0<\lambda$, $\mu\leq 100$. The optimum solution to this particular problem is $\lambda=1.1$, $\mu=0.1$, and $A_q=1$.

Frugivores may travel far away from the plants and hence will disperse the seeds far and wide. This feeding behaviour typically follows a Lévy distribution [43–45]. In the following we present some basic facts about it.

2.1. Lévy Distribution. The Lévy distribution is a probability density distribution for random variables. Here the random variables represent the directions of flights of arbitrary birds. This function ranges over real numbers in the domain represented by the problem search space.

The flight lengths of the agents served by the plants follow a heavy tailed power law distribution [14], represented by

$$L(s) \sim |s|^{-1-\beta},\tag{5}$$

where L(s) denotes the Lévy distribution with index $\beta \in (0,2)$. Lévy flights are unique arbitrary excursions whose step lengths are drawn from (5). An alternative form of Lévy distribution is [14]

$$L(s, \gamma, \mu) = \begin{cases} \sqrt{\frac{\gamma}{2\pi}} \left(\frac{1}{(s-\mu)}\right)^{3/2} \\ \cdot \exp\left[-\frac{\gamma}{2(s-\mu)}\right], & 0 < \mu < s < \infty, \end{cases}$$

$$0, \qquad \text{Otherwise.}$$

This implies that

$$\lim_{s \to \infty} L(s, \gamma, \mu) \approx \sqrt{\frac{\gamma}{2\pi}} \left(\frac{1}{s}\right)^{3/2}.$$
 (7)

In terms of the Fourier transform [14], the limiting value of L(s) can be written as

$$\lim_{s \to \infty} L(s) = \frac{\alpha \beta \Gamma(\beta) \sin(\pi \beta/2)}{\pi |s|^{1+\beta}},$$
 (8)

where $\Gamma(\beta)$ is the Gamma function [46], defined by

$$\Gamma(\beta) = \int_0^\infty x^{\beta - 1} e^{-x} dx. \tag{9}$$

The steps L(s) are generated by Mantegna's algorithm [14]. This algorithm ensures that the behaviour of Lévy flights is symmetric and stable as shown in Figure 3(b).

3. Strawberry Plant Propagation Algorithm: The Feeding Station Model

We assume that the arrival of different agents (birds and animals) to the plants to feed is according to the Poisson distribution [40]. As per the solution of problem (4), the mean arrival rate is $\lambda = 1.1$, and NP = 10 is the size of the agents population. Let k = 1, 2, ..., A be the possible numbers of agents visiting the plants per unit time. With these assumptions the graphic representation of (1) results in Figure 2.

As already stated, it is essential in this algorithm to balance exploration and exploitation. To this end, we choose a threshold value of the Poisson probability that dictates how much exploration and exploitation are done during the search. The probability $Poiss(\lambda) < 0.05$ means that

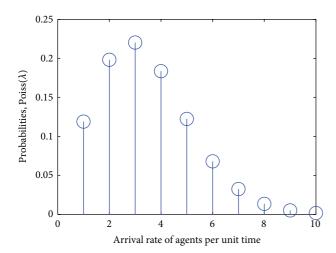


FIGURE 2: Distribution of agents arriving at strawberry plants to eat fruit and disperse seeds.

exploitation is covered. In this case, (10) below is used, which helps the algorithm to search locally:

$$x_{i,j}^* = \begin{cases} x_{i,j} + \xi_j \left(x_{i,j} - x_{l,j} \right) & \text{if PR} \le 0.8; \ j = 1, 2, \dots, n; \\ \\ i, l = 1, 2, \dots, \text{NP}; \ i \ne l, \\ \\ x_{i,j}, & \text{Otherwise,} \end{cases}$$
(10)

where PR denotes the rate of dispersion of the seeds locally, around SP; $x_{i,j}^*$ and $x_{i,j} \in [a_j \ b_j]$ are the jth coordinates of the seeds X_i^* and X_i , respectively; a_j and b_j are the jth lower and upper bounds defining the search space of the problem and $\xi_i \in [-1 \ 1]$. The indices l and i are mutually exclusive.

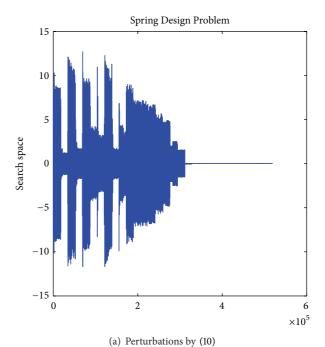
On the other hand, if $Poiss(\lambda) \ge 0.05$ then global dispersion of seeds becomes more prominent. This is implemented by using the following equation:

$$x_{i,j}^* = \begin{cases} x_{i,j} + L_i \left(x_{i,j} - \theta_j \right) & \text{if PR} \le 0.8, \ \theta_j \in \left[a_j \ b_j \right], \\ & i = 1, 2, \dots, \text{NP}; \\ & j = 1, 2, \dots, n, \\ x_{i,j}, & \text{Otherwise,} \end{cases}$$
(11)

where L_i is a step drawn from the Lévy distribution [14] and θ_j is a random coordinate within the search space. Equations (10) and (11) perturb the current solution, the results of which can be seen in Figures 3(a) and 3(b), respectively.

As mentioned in Algorithm 1, we first collect the best solutions from the first NP trial runs to form a population of potentially good solutions denoted by pop_{best} . The convergence rate of SbPPA is shown in Figures 4 and 5 for different test problems used in our experiments (see Appendices). The statistics values best, worst, mean, and standard deviation are calculated based on pop_{best} .

The seed-based propagation process of SP can be represented in the following steps.



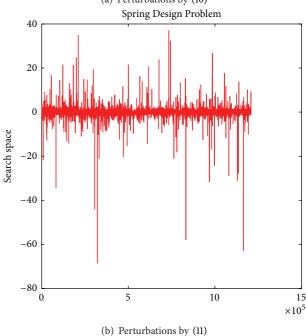


FIGURE 3: Overall performance of SbPPA on Spring Design Problem.

- (1) The dispersal of seeds in the neighbourhood of the SP, as shown in Figure 1(e), is carried out either by fruits fallen from strawberry plants after they become ripe or by agents. The step lengths for this phase are calculated using (10).
- (2) Seeds are spread globally through agents, as shown in Figures 1(c) and 1(d). The step lengths for these travelling agents are drawn from the Lévy distribution [14].
- (3) The probabilities, $Poiss(\lambda)$, that a certain number k of agents will arrive to SP to eat fruits and disperse it,

```
(1) NP \leftarrow Population size, r \leftarrow Counter of trial runs, MaxExp \leftarrow Maximum experiments
(2) for r = 1: MaxExp do
     if r \leq NP then
         Create a random population of seeds pop = \{X_i \mid i = 1, 2, ..., NP\},
(4)
         using (12) and collect the best solutions from each trial run, in pop<sub>best</sub>.
(5)
         Evaluate the population pop.
(6)
       end if
      while r > NP do
(7)
         Use updated population popbest.
(8)
(9)
      end while
(10)
       while (the stopping criteria is not satisfied) do
(11)
         for i = 1 to NP do
(12)
            if Poiss(\lambda)_i \ge 0.05, then,
                                                             ▷ (Global or local seed dispersion)
               for i = 1 to n do
(13)
                                                                  \triangleright (n is number of dimensions)
                                                                        ▷ (PR = Perturbation Rate)
(14)
                  if rand \leq PR then,
(15)
                    Update the current entry according to (11)
(16)
                  end if
(17)
               end for
(18)
             else
(19)
               for j = 1 to n do
(20)
                  if rand \leq PR then,
(21)
                     Update the current entry according to (10)
(22)
                   end if
(23)
               end for
(24)
             end if
(25)
          end for
(26)
          Update current best
(27)
       end while
       Return: Updated population and global best solution.
(28)
(29) end for
```

ALGORITHM 1: Seed-based Plant Propagation Algorithm (SbPPA) [47].

is used as a balancing factor between exploration and exploitation.

For implementation purposes, we assume that each SP produces one fruit, and each fruit is assumed to have one seed; by a solution X_i we mean the current position of the ith seed to be dispersed. The number of seeds in the population is denoted by NP. Initially we generate a random population of NP seeds using

$$x_{i,j} = a_j + (b_j - a_j) \eta_j, \quad j = 1, ..., n,$$
 (12)

where $x_{i,j} \in [a_j \ b_j]$ is the jth coordinate of solution X_i , a_j and b_j are the jth coordinates of the bounds describing the search space of the problem, and $\eta_j \in (0 \ 1)$. This means that $X_i = [x_{i,j}]$, for $j = 1, \ldots, n$, represents the position of the jth seed in population pop.

4. Experimental Settings and Discussion

In our experiments we tested SbPPA against some recently developed algorithms and some well established and standard ones. Our set of test problems includes benchmark constrained and unconstrained optimization problems [22, 30, 48, 49]. The results are compared in terms of statistics (best, worst, mean and standard deviation) for solutions obtained

by SbPPA; ABC [18, 50]; PSO [51]; FF [21]; HPA [29]; SSO-C [22]; Classical Evolutionary Programming (CEP) [30]; and Fast Evolutionary Programming (FEP) [30]. The detailed descriptions of these problems are given in Appendices.

In Tables 4 and 7, the significance of results is shown in terms of win/tie/loss (see Table 2 in [52]) according to the following notations:

- (i) (+) when SbPPA is better;
- (ii) (≈) when the results are approximately the same as those obtained with SbPPA;
- (iii) (-) when SbPPA is worse.

Moreover, in Tables 5 and 6 the significance of results obtained with SbPPA is highlighted.

4.1. Parameter Settings. The parameter settings are given in Tables 1-3.

5. Conclusion

In this paper, a new metaheuristic referred to as the Seed-Based Plant Propagation Algorithm (SbPPA) [47] has been proposed. Plants have evolved a variety of ways to propagate. Propagation through seeds is perhaps the most common of them all and one which takes advantage of all sorts of

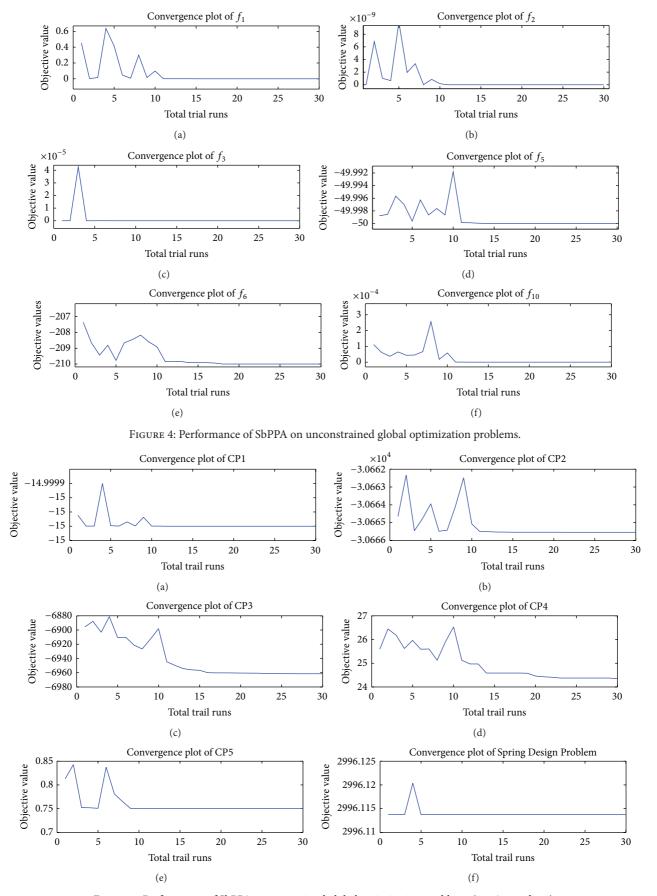


FIGURE 5: Performance of SbPPA on constrained global optimization problems (see Appendices).

TABLE 1: Parameters used for each algorithm for solving unconstrained global optimization problems $f_1 - f_{10}$. All experiments are repeated 30 times.

BSO [15 20]	00 81 Jav	LIPA (20)	Shph [47]
F5O [16, 29]	ABC[18, 29]	пРА [29]	50FFA [47]
M = 100	SN = 100	Agents = 100	NP = 10
G (Dimension \times 20, 000)	$M_{\text{CN}} = (\text{Dimension} \times 20,000)$	The state of $\frac{\text{(Dimension} \times 20,000)}{\text{(Dimension} \times 20,000)}$	Iteration number – (Dimension $\times 20,000$)
C _{max} – M	SN	Agents	NP
$c_1 = 2$	MR = 0.8	$c_1 = 2$	PR = 0.8
•	(SN × dimension)	•	
$c_2 = 2$	Limit = $\frac{1}{2}$	$c_2 = 2$	$Poiss(\lambda) = 0.05$
:	7	(1	
$_{147}$ (G_{max} – iteration _{index})		$G_{1}:=$ (Agents × dimension)	
	I	DIIIII =	I
Cmax		1	
1		(Iteration number – iteration _{index})	
	I	W =Iteration number	I
		Itelation manner	

Table 2: Experimental setup used for each algorithm for solving unconstrained global optimization problems f_{11} – f_{26} . All experiments are repeated 50 times.

CEP [30, 53, 54]	FEP [30]	SbPPA
Population size $\mu = 100$	Population size $\mu = 100$	NP = 10
Tournament size $q = 10$	Tournament size $q = 10$	PR = 0.8
$\eta = 3.0$	$\eta = 3.0$	$Poiss(\lambda) = 0.05$

agents ranging from wind to water, birds, and animals. The strawberry plant uses both runners and seeds to propagate. Here we consider the propagation through seeds that the strawberry plant has evolved, to design an efficient optimization algorithm.

To capture the dispersal process, we adopt a queuing approach which, given the extent of fruit produced, indicates the extent of seeds dispersed and hence the effectiveness of the search/optimization algorithm based on this process. Looking at the random process of agents using the plants (feeding station) it is reasonable to assume that it is of the Poisson type. On the other hand, the time taken by agents in successfully eating fruit and leaving to disperse its seeds, in other words the service time for agents, is expressed by a random variable which follows the exponential probability distribution. To this end, we choose a threshold value of the Poisson probability that dictates how much exploration and exploitation are done during the search. An alternative strategy has been adopted here. This strategy consists in making sure that the initial population is as good as the user can afford it to be by using best solutions found so far. The effects of this strategy on convergence are shown through convergence plots of Figures 4 and 5, for some of the solved problems. SbPPA is easy to implement as it requires less arbitrary parameter settings than other algorithms. The success rate of SbPPA increases as it gets its population of best solutions. It has been implemented for both unconstrained and constrained optimization problems. Its performance, compared to that of other algorithms, points to SbPPA as being superior.

Appendices

A. Unconstrained Global Optimization Problems

See Tables 8 and 9.

B. Set of Constrained Global Optimization Problems Used in Our Experiments

B.1. CP1. Consider the following:

Min
$$f(x) = 5\sum_{d=1}^{4} x_d - 5\sum_{d=1}^{4} x_d^2 - \sum_{d=5}^{13} x_d$$
,

subject to
$$g_{1}(x) = 2x_{1} + 2x_{2} + x_{10} + x_{11} - 10 \le 0,$$

$$g_{2}(x) = 2x_{1} + 2x_{3} + x_{10} + x_{12} - 10 \le 0,$$

$$g_{3}(x) = 2x_{2} + 2x_{3} + x_{11} + x_{12} - 10 \le 0,$$

$$g_{4}(x) = -8x_{1} + x_{10} \le 0,$$

$$g_{5}(x) = -8x_{2} + x_{11} \le 0,$$

$$g_{6}(x) = -8x_{3} + x_{12} \le 0,$$

$$g_{7}(x) = -2x_{4} - x_{5} + x_{10} \le 0,$$

$$g_{8}(x) = -2x_{6} - x_{7} + x_{11} \le 0,$$

$$g_{9}(x) = -2x_{8} - x_{9} + x_{12} \le 0,$$
(B.1)

where bounds are $0 \le x_i \le 1$ (i = 1, ..., 9, 13), $0 \le x_i \le 100$ (i = 10, 11, 12). The global optimum is at $x^* = (1, 1, 1, 1, 1, 1, 1, 1, 1, 3, 3, 3, 1)$, $f(x^*) = -15$.

B.2. CP2. Consider the following:

Min
$$f(x) = 5.3578547x_2 + 0.8356891x_1x_5 \\ + 37.293239x_1 - 40792.141,$$
 subject to
$$g_1(x) = 85.334407 + 0.0056858x_2x_5 \\ + 0.0006262x_1x_4 - 0.0022053x_3x_5 \\ - 92 \le 0,$$

$$g_2(x) = -85.334407 - 0.0056858x_2x_5 \\ - 0.0006262x_1x_4 + 0.0022053x_3x_5 \\ \le 0,$$

$$g_3(x) = 80.51249 + 0.0071317x_2x_5 \\ + 0.0029955x_1x_2 - 0.0021813x_2 \\ - 110 \le 0,$$

$$g_4(x) = -80.51249 - 0.0071317x_2x_5 \\ + 0.0029955x_1x_2 - 0.0021813x_2 \\ + 90 \le 0,$$

$$g_5(x) = 9.300961 - 0.0047026x_3x_5 \\ - 0.0012547x_1x_3 - 0.0019085x_3x_4 \\ - 25 \le 0,$$

$$g_6(x) = -9.300961 - 0.0047026x_3x_5 \\ - 0.0012547x_1x_3 - 0.0019085x_3x_4 \\ + 20 \le 0,$$
 (B.2)

Table 3: Parameters used for each algorithm for solving constrained optimization problems. All experiments are repeated 30 times.

PSO [51]	ABC [50]	FF [21]	SSO-C [22]	SbPPA [47]
M = 250	SN = 40	Fireflies = 25	N = 50	NP = 10
$G_{\text{max}} = 300$	MCN = 6000	Iteration number = 2000	Iteration number = 500	Iteration number = 800
$c_1 = 2$	MR = 0.8	q = 1.5	PF = 0.7	PR = 0.8
$c_2 = 2$	_	$\alpha = 0.001$	_	$Poiss(\lambda) = 0.05$
Weight factors = 0.9 to 0.4	_	_	_	_

TABLE 4: Results obtained by SbPPA, HPA, PSO, and ABC. All problems in this table are unconstrained.

Fun.	Dim	Algorithm	Best	Worst	Mean	SD
		ABC	(+) 0.0129	(+) 0.6106	(+) 0.1157	(+) 0.111
£	4	PSO	(-) 6.8991E - 08	(+) 0.0045	(+) 0.001	(+) 0.0013
f_1	4	HPA	(+) 2.0323E - 06	(+) 0.0456	(+) 0.009	(+) 0.0122
		SbPPA	1.08E - 07	7.05E - 06	3.05E - 06	3.14E - 06
		ABC	(+) 1.2452E - 08	(+) 8.4415 <i>E</i> – 06	(+) 1.8978 <i>E</i> – 06	(+) 1.8537E - 06
f_2	2	PSO	(≈) 0	(≈) 0	(≈) 0	(≈)0
	2	HPA	(≈) 0	(≈) 0	(≈) 0	(≈) 0
		SbPPA	0	0	0	0
		ABC	(≈) 0	(+) 4.8555 <i>E</i> – 06	(+) 4.1307E - 07	(+) 1.2260E - 06
C	2	PSO	(≈) 0	(+) 3.5733E - 07	(+) 1.1911E - 08	(+) 6.4142E - 08
f_3	2	HPA	(≈) 0	(≈) 0	(≈) 0	(≈) 0
		SbPPA	0	0	0	0
		ABC	(≈) −1.03163	(≈) −1.03163	(≈) −1.03163	(≈)0
<u>r</u>	2	PSO	(≈) −1.03163	(≈) −1.03163	(≈) −1.03163	(≈)0
f_4	2	HPA	(≈) −1.03163	(≈) −1.03163	(≈) −1.03163	(≈)0
		SbPPA	-1.031628	-1.031628	-1.031628	0
		ABC	(≈) −50.0000	(≈) −50.0000	(≈) −50.0000	(-)0
f_5	6	PSO	(≈) −50.0000	(≈) −50.0000	$(\approx) -50.0000$	(-) 0
	6	HPA	(≈) −50.0000	(≈) −50.0000	$(\approx) -50.0000$	(-) 0
		SbPPA	-50.0000	-50.0000	-50.0000	5.88E - 09
		ABC	(+) -209.9929	(+) -209.8437	(+) -209.9471	(+) 0.044
ſ	10	PSO	(≈) −210.0000	$(\approx) -210.0000$	$(\approx) -210.0000$	(-) 0
f_6	10	HPA	(≈) −210.0000	$(\approx) -210.0000$	$(\approx) -210.0000$	(+)1
		SbPPA	-210.0000	-210.0000	-210.0000	4.86E - 06
		ABC	(+) 2.6055 <i>E</i> – 16	(+) 5.5392 <i>E</i> – 16	(+) 4.7403 <i>E</i> –16	(+) 9.2969 <i>E</i> – 17
f_7	30	PSO	(≈) 0	(≈) 0	(≈) 0	(≈) 0
J_7	30	HPA	(≈) 0	(≈) 0	(≈) 0	(≈) 0
		SbPPA	0	0	0	0
		ABC	(+) 2.9407E - 16	(+) 5.5463E - 16	(+) 4.8909E - 16	(+) 9.0442 <i>E</i> – 17
£	30	PSO	(≈) 0	(≈) 0	(≈) 0	(≈) 0
f_8	30	HPA	(≈) 0	(≈) 0	(≈) 0	(≈) 0
		SbPPA	0	0	0	0
		ABC	(≈) 0	(+) 1.1102E - 16	(+) 9.2519E - 17	(+) 4.1376E - 17
£	30	PSO	(≈) 0	(+) 1.1765E - 01	(+) 2.0633E - 02	(+) 2.3206E - 02
f_9	30	HPA	(≈) 0	(≈) 0	(≈) 0	(≈) 0
		SbPPA	0	0	0	0
		ABC	(+) 2.9310 <i>E</i> – 14	(+) 3.9968 <i>E</i> – 14	(+) 3.2744 <i>E</i> – 14	(+) 2.5094 <i>E</i> – 15
ſ	30	PSO	$(\approx) 7.9936E - 15$	(+) 1.5099E - 14	(+) 8.5857E - 15	(+) 1.8536 <i>E</i> – 15
f_{10}	30	HPA	$(\approx) 7.9936E - 15$	(+) 1.5099E - 14	(+) 1.1309E - 14	(+) 3.54E - 15
		SbPPA	7.994E - 15	7.99361E - 15	7.994E - 15	7.99361 <i>E</i> – 15

TABLE 5: Results obtained by SbPPA, CEP, and FEP. All problems in this table are unconstrained [30].

Function number	Algorithm	Maximum generations	Mean	SD
	CEP		2.60E - 03	1.70E - 04
f_{11}	FEP	2000	8.10E - 03	7.70E - 04
	SbPPA		9.45E - 13	4.08E - 12
	CEP		2	1.2
f_{12}	FEP	5000	0.3	0.5
	SbPPA		3.93E - 02	3.76E - 02
	CEP		6.17	13.61
f_{13}	FEP	20000	5.06	5.87
	SbPPA		1.86E + 01	2.25E + 00
	CEP		577.76	1125.76
f_{14}	FEP	1500	0	0
	SbPPA		0	0
	CEP		1.80E - 02	6.40E - 03
f_{15}	FEP	3000	7.60E - 03	2.60E - 03
	SbPPA		3.61E - 03	1.31E - 03
	CEP		-7.92E + 03	6.35E + 02
f_{16}	FEP	9000	-1.26E + 04	5.26E + 01
	SbPPA		-1.16E + 04	6.04E + 01
	CEP		89	23.1
f_{17}	FEP	5000	4.60E - 02	1.20E - 02
	SbPPA		8.73E + 00	9.88E - 01
	CEP		0.398	1.50E - 07
f_{18}	FEP	100	0.398	1.50E - 07
	SbPPA		3.98E - 01	0

Table 6: Results obtained by SbPPA, CEP, and FEP. All problems in this table are unconstrained [30].

Function number	Algorithm	Maximum generations	Mean	SD
	CEP		3	0
f_{19}	FEP	100	3.02	0.11
	SbPPA		3	3.05E - 15
	CEP		-3.86E + 00	1.40E - 02
f_{20}	FEP	100	-3.86E + 00	1.40E - 05
	SbPPA		-3.86E + 00	2.75E - 15
	CEP		-3.28E + 00	5.80E - 02
f_{21}	FEP	200	-3.27E + 00	5.90E - 02
	SbPPA		-3.32E + 00	2.91E - 14
	CEP		-6.86E + 00	2.67E + 00
f_{22}	FEP	100	-5.52E + 00	1.59E + 00
	SbPPA		-1.02E + 01	4.30E - 09
	CEP		-8.27	2.95
f_{23}	FEP	100	-5.52	2.12
	SbPPA		-1.04E + 01	7.73E - 09
	CEP		-9.1	2.92
f_{24}	FEP	100	-6.57	3.14
	SbPPA		-1.05E + 01	1.03E - 07
	CEP		1.66	1.19
f_{25}	FEP	100	1.22	0.56
	SbPPA		9.98E - 01	1.13E - 16
	CEP		4.70E - 04	3.00E - 04
f_{26}	FEP	4000	5.00E - 04	3.20E - 04
	SbPPA		3.07E - 04	6.80E - 15

Table 7: Results obtained by SbPPA, PSO, ABC, FF, and SSO-C. All problems in this table are standard constrained optimization problems.

Fun. name	O-4:1	Algorithm	Best	Mean	Worst	SD
Tun. name	Optimal	PSO	(≈) −15	(≈) -15	(≈) −15	(-) 0
		ABC	(≈) −15 (≈) −15	(≈) −15 (≈) −15	(≈) −15 (≈) −15	(-) 0
CP1	-15	FF	(+) 14.999	(+) 14.988	(~) -13 (+) 14.798	(-) 0 $(+)$ 6.40 E – 07
		SSO-C	(≈) −15	(≈) −15	(≈) −15	(+) 0.40L - 07 (-) 0
		SbPPA	(~) -13 -15	(~) -13 -15	(~) -13 -15	1.95E - 15
		PSO	(≈) −30665.5	(+) -30662.8	(+) -30650.4	(+) 5.20E - 02
		ABC	(≈) −30665.5	(+) -30662.8 (+) -30664.9	(+) -30659.1	(+) 3.20 E - 02 $(+)$ 8.20 E - 02
CP2	-30665.539	FF	$(\approx) -3.07E + 04$	(+) -30662	(+) -30649	(+) 5.20E - 02 (+) 5.20E - 02
		SSO-C	$(\approx) -3.07E + 04$ $(\approx) -3.07E + 04$	(≈) −30665.5	(+) -30665.1	(+) 3.20E - 02 (+) 1.10E - 04
		SbPPA	-30665.5	-30665.5	-30665.5	(+) 1.10E - 04 2.21E - 06
		PSO				
		ABC	(+) $-6.96E + 03$ $(-)$ -6961.81	(+) -6958.37 (+) 6958.02	(+) -6942.09 (+) -6955.34	(+) 6.70E - 02 (-) 2.10E - 02
CP3	-6961.814	FF	(+) -6959.99	(+) -6958.02 $(+)$ $-6.95E + 03$	(+) -6933.34 (+) -6947.63	(-) 2.10E $-$ 02 (-) 3.80E $-$ 02
		SSO-C				
		SbPPA	(-) -6961.81	(+) -6961.01	(+) -6960.92	(-) 1.10E - 03
			-6961.5	-6961.38	-6961.45	0.043637
		PSO	(-) 24.327	(+) 2.45E + 01	(+) 24.843	(+) 1.32E - 01
CP4	24.306	ABC	(+) 24.48	(+) 2.66E + 01	(+) 28.4	(+) 1.14
	24.500	FF SSO C	(-) 23.97	(+) 28.54	(+) 30.14	(+) 2.25
		SSO-C	(-) 24.306	(-) 24.306	(-) 24.306	(-) 4.95 E – 05
		SbPPA	24.34442	24.37536	24.37021	0.012632
		PSO	(≈) −0.7499	(+) -0.749	(+) -0.7486	(+) 1.20E - 03
CP5	-0.7499	ABC	(≈) −0.7499	(+) -0.7495	(+) -0.749	(+) 1.67E - 03
	-0.7499	FF	(+) -0.7497	(+) -0.7491	(+) -0.7479	(+) 1.50E - 03
		SSO-C	(≈) −0.7499	(≈) −0.7499	(≈) −0.7499	(-) 4.10E - 09
		SbPPA	0.7499	0.749901	0.7499	1.66E - 07
		PSO	(+) 0.012858	(+) 0.014863	(+) 0.019145	(+) 0.001262
Spring Design Problem	Not known	ABC	(≈) 0.012665	(+) 0.012851	(+) 0.01321	(+) 0.000118
opinig Beoign Froblem	Not kilowii	FF	(≈) 0.012665	(+) 0.012931	(+) 0.01342	(+) 0.001454
		SSO-C	(≈) 0.012665	(+) 0.012765	(+) 0.012868	(+) 9.29 E – 05
		SbPPA	0.012665	0.012666	0.012666	3.39E - 10
		PSO	(+) 1.846408	(+) 2.011146	(+) 2.237389	(+) 0.108513
Welded beam design problem	Not len oven	ABC	(+) 1.798173	(+) 2.167358	(+) 2.887044	(+) 0.254266
weided beam design problem	Not known	FF	(+) 1.724854	(+) 2.197401	(+) 2.931001	(+) 0.195264
		SSO-C	(≈) 1.724852	(+) 1.746462	(+) 1.799332	(+) 0.02573
		SbPPA	1.724852	1.724852	1.724852	4.06E - 08
		PSO	(+) 3044.453	(+) 3079.262	(+) 3177.515	(+) 26.21731
Speed reducer design antimication	NI (I	ABC	(+) 2996.116	(+) 2998.063	(+) 3002.756	(+) 6.354562
Speed reducer design optimization	Not known	FF	(+) 2996.947	(+) 3000.005	(+) 3005.836	(+) 8.356535
		SSO-C	(≈) 2996.113	(≈) 2996.113	(≈) 2996.113	(+) 1.34E - 12
		SbPPA	2996.114	2996.114	2996.114	0

Fun.	Fun. name	D	C	Range	Min	Formulation
f_1	Colville	4	UN	$[-1010]^D$	0	$f(x) = 100(x_1^2 - x_2) + (x_1 - 1)^2 + (x_3 - 1)^2 + 90(x_3^2 - x_4)^2 + 10.1((x_2 - 1)^2 + (x_4 - 1)^2) + 19.8(x_2 - 1)(x_4 - 1)$
$\overline{f_2}$	Matyas	2	UN	$[-1010]^D$	0	$f(x) = 0.26(x_1^2 + x_2^2) - 0.48x_1x_2$
f_3	Schaffer	2	MN	$[-100100]^D$	0	$f(x) = 0.5 + \frac{\sin^2\left(\sqrt{\sum_{i=1}^n x_i^2}\right) - 0.5}{\left(1 + 0.001\left(\sum_{i=1}^n x_i^2\right)\right)^2}$
f_4	Six Hump Camel Back	2	MN	$[-55]^{D}$	-1.03163	$f(x) = 4x_1^2 - 2.1x_1^4 + \frac{1}{3}x_1^6 + x_1x_2 - 4x_2^2 + 4x_2^4$
f_5	Trid6	6	UN	$[-3636]^D$	-50	$f(x) = \sum_{i=1}^{6} (x_i - 1)^2 - \sum_{i=2}^{6} x_i x_{i-1}$
f_6	Trid10	10	UN	$[-100100]^D$	-210	$f(x) = \sum_{i=1}^{10} (x_i - 1)^2 - \sum_{i=2}^{10} x_i x_{i-1}$
f_7	Sphere	30	US	$[-100\ 100]^D$	0	$f(x) = \sum_{i=1}^{n} x_i^2$
f_8	SumSquares	30	US	$[-1010]^D$	0	$f(x) = \sum_{i=1}^{n} i x_i^2$
f_9	Griewank	30	MN	$[-600600]^D$	0	$f(x) = \frac{1}{4000} \sum_{i=1}^{n} x_i^2 - \prod_{i=1}^{n} \cos\left(\frac{x_i}{\sqrt{i}}\right) + 1$
f_{10}	Ackley	30	MN	$[-32\ 32]^D$	0	$f(x) = -20 \exp\left(-0.2\sqrt{\frac{1}{n}\sum_{i=1}^{n}x_i^2}\right) - \exp\left(\frac{1}{n}\sum_{i=1}^{n}\cos\left(2\pi x_i\right)\right) + 20 + e$

Table 8: Unconstrained global optimization problems (Set-1) used in our experiments.

where $78 \le x_1 \le 102$, $33 \le x_2 \le 45$, and $27 \le x_i \le 45$ (i = 3, 4, 5). The optimum solution is $x^* = (78, 33, 29.995256025682, 45, 36.775812905788)$, where $f(x^*) = -30665.539$. Constraints g_1 and g_6 are active.

B.3. CP3. Consider the following:

Min
$$f(x) = (x_1 - 10)^3 + (x_2 - 20)^3$$
,
subject to $g_1(x) = -(x_1 - 5)^2 - (x_2 - 5)^2 + 100 \le 0$,
 $g_2(x) = (x_1 - 6)^2 + (x^2 - 5)^2 - 82.81 \le 0$,
(B.3)

where $13 \le x_1 \le 100$ and $0 \le x_2 \le 100$. The optimum solution is $x^* = (14.095, 0.84296)$ where $f(x^*) = -6961.81388$. Both constraints are active.

B.4. CP4. Consider the following:

Min
$$f(x) = x_1^2 + x_2^2 + x_1 x_2 - 14x_1 - 16x_2$$

$$+ (x_3 - 10)^2 + 4(x_4 - 5)^2 + (x_5 - 3)^2$$

$$+ 2(x_6 - 1)^2 + 5x_7^2 + 7(x_8 - 11)^2$$

$$+ 2(x_9 - 10)^2 + (x_{10} - 7)^2 + 45,$$
subject to
$$g_1(x) = -105 + 4x_1 + 5x_2 - 3x_7 + 9x_8 \le 0,$$

$$g_2(x) = 10x_1 - 8x_2 - 17x_7 + 2x_8 \le 0,$$

$$g_{3}(x) = -8x_{1} + 2x_{2} + 5x_{9} - 2x_{10} - 12 \le 0,$$

$$g_{4}(x) = 3(x_{1} - 2)^{2} + 4(x_{2} - 3)^{2} + 2x_{3}^{2}$$

$$-7x_{4} - 120 \le 0,$$

$$g_{5}(x) = 5x_{1}^{2} + 8x_{2} + (x_{3} - 6)^{2} - 2x^{4}$$

$$-40 \le 0,$$

$$g_{6}(x) = x_{1}^{2} + 2(x_{2} - 2)^{2} - 2x_{1}x_{2} + 14x_{5}$$

$$-6x_{6} \le 0,$$

$$g_{7}(x) = 0.5(x_{1} - 8)^{2} + 2(x_{2} - 4)^{2} + 3x_{5}^{2}$$

$$-x_{6} - 30 \le 0,$$

$$g_{8}(x) = -3x_{1} + 6x_{2} + 12(x_{9} - 8)^{2} - 7x_{10} \le 0,$$

where $-10 \le x_i \le 10$ (i = 1, ..., 10). The global optimum is $x^* = (2.171996, 2.363683, 8.773926, 5.095984, 0.9906548, 1.430574, 1.321644, 9.828726, 8.280092, 8.375927), where <math>f(x^*) = 24.3062091$. Constraints g_1, g_2, g_3, g_4, g_5 , and g_6 are active.

B.5. CP5. Consider the following:

Min
$$f(x) = x_1^2 + (x_2 - 1)^2$$
,
subject to $g_1(x) = x_2 - x_1^2 = 0$, (B.5)

where $1 \le x_1 \le 1$, $1 \le x_2 \le 1$. The optimum solution is $x^* = (\pm 1/\sqrt{(2)}, 1/2)$, where $f(x^*) = 0.7499$.

Table 9: Unconstrained global optimization problems (Set-2) used in our experiments [30].

Fun. number	Range	D	Function	Formulation	$f_{ m min}$
f_{11}	$[-10, 10]^D$	30	Schwefel's Problem 2.22	$f(x) = \sum_{i=1}^{n} x_i + \prod_{i=1}^{n} x_i $	0
f_{12}	$[-100, 100]^D$	30	Schwefel's Problem 2.21	$f(x) = \max_{i} \{ x_i , 1 \le i \le n\}$	0
f_{13}	$[-10, 10]^D$	30	Rosenbrock	$f(x) = \sum_{i=1}^{n-1} \left[100 \left(x_{i+1} - x_i^2 \right)^2 + \left(x_i - 1 \right)^2 \right]$	0
f_{14}	$[-100, 100]^D$	30	Step	$f(x) = \sum_{i=1}^{n} (\lfloor x_i + 0.5 \rfloor)^2$	0
f_{15}	$[-1.28, 1.28]^D$	30	Quartic (noise)	$f(x) = \sum_{i=1}^{n} ix_i^4 + \text{random}[0, 1)$	0
f_{16}	$[-500, 500]^D$	30	Schwefel	$f(x) = -\sum_{i=1}^{n} x_i \sin\left(\sqrt{ x_i }\right)$	-12569.5
f_{17}	$[-5.12, 5.12]^D$	30	Rastrigin	$f(x) = [x_i^2 - 10\cos(2\pi x_i) + 10]$	0
f_{18}	$[-5, 10] \times [0, 15]$	2	Branin	$f(x) = \left(x_2 - \frac{5.1}{4\pi^2}x_1^2 + \frac{5}{\pi}x_1 - 6\right)^2 + 10\left(1 - \frac{1}{8\pi}\right)\cos x_1 + 10$	0.398
f_{19}	$[-2, 2]^D$	2	Goldstein- Price	$f(x) = \left[1 + (x_1 + x_2 + 1)^2 \left(19 - 14x_1 + 3x_1^2 - 14x_2 + 6x_1x_2 + 3x_2^2\right)\right]$ $\times \left[30 + \left(2x_1 - 3x_2\right)^2 \left(18 - 32x_1 + 12x_1^2 + 48x_2 - 36x_1x_2 + 27x_2^2\right)\right]$	3
f_{20}	$[0,1]^{D}$	4	Hartman's Family $(n = 3)$	$f(x) = -\sum_{i=1}^{4} c_i \exp\left[\sum_{j=1}^{3} a_{ij} (x_j - p_{ij})^2\right]$	-3.86
f_{21}	$[0,1]^D$	6	Hartman's Family $(n = 6)$	$f(x) = -\sum_{i=1}^{4} c_i \exp\left[\sum_{j=1}^{6} a_{ij} (x_j - p_{ij})^2\right]$	-3.32
f_{22}	$[0,10]^D$	4	Shekel's Family (<i>m</i> = 5)	$f(x) = -\sum_{i=1}^{5} \left[(x - a_i)(x - a_i)^T + c_i \right]^{-1}$	-10
f_{23}	$[0,10]^D$	4	Shekel's Family (<i>m</i> = 7)	$f(x) = -\sum_{i=1}^{7} [(x - a_i)(x - a_i)^T + c_i]^{-1}$	-10
f_{24}	$[0,10]^D$	4	Shekel's Family (<i>m</i> = 10)	$f(x) = -\sum_{i=1}^{10} \left[(x - a_i) (x - a_i)^T + c_i \right]^{-1}$	-10
f_{25}	[-65.536, 65.536] ^D	2	Shekel's Foxholes	$f(x) = \left[\frac{1}{500} + \sum_{j=1}^{25} \frac{1}{j + \sum_{i=1}^{2} (x_i - a_{ij})^6}\right]^{-1}$	1
f_{26}	$[-5, 5]^D$	4	Kowalik	$f(x) = \sum_{i=1}^{11} \left[a_i - \frac{x_1 (b_i^2 + b_i x_2)}{b_i^2 + b_i x_3 + x_4} \right]^2$	0.0003075

B.6. Welded Beam Design Optimisation. The welded beam design is a standard test problem for constrained design optimisation [55, 56]. There are four design variables: the width w and length L of the welded area and the depth d and thickness h of the main beam. The objective is to minimise the overall fabrication cost, under the appropriate constraints of shear stress τ , bending stress σ , buckling load P, and maximum end deflection δ . The optimization model is summarized as follows, where $x^T = (w, L, d, h)$:

Minimise
$$f(x) = 1.10471w^2L + 0.04811dh(14.0 + L)$$
,
subject to $g_1(x) = w - h \le 0$,

$$\begin{split} g_2\left(x\right) &= \delta\left(x\right) - 0.25 \leq 0, \\ g_3\left(x\right) &= \tau\left(x\right) - 13,600 \leq 0, \\ g_4\left(x\right) &= \sigma\left(x\right) - 30,000 \leq 0, \\ g_5\left(x\right) &= 1.10471w^2 + 0.04811dh\left(14.0 + L\right) \\ &\quad - 5.0 \leq 0, \\ g_6\left(x\right) &= 0.125 - w \leq 0, \\ g_7\left(x\right) &= 6000 - P\left(x\right) \leq 0, \end{split} \tag{B.6}$$

where

$$\sigma(x) = \frac{504,000}{hd^2}, \qquad D = \frac{1}{2}\sqrt{L^2 + (w+d)^2},$$

$$Q = 6000\left(14 + \frac{L}{2}\right), \qquad \delta = \frac{65,856}{30,000hd^3},$$

$$J = \sqrt{2}wL\left(\frac{L^2}{6} + \frac{(w+d)^2}{2}\right), \qquad \alpha = \frac{6000}{\sqrt{2}wL},$$

$$\beta = \frac{QD}{J}, \qquad P = 0.61423 \times 10^6 \frac{dh^3}{6} \left(1 - \frac{\sqrt[d]{30/48}}{28}\right),$$

$$\tau(x) = \sqrt{\alpha^2 + \frac{\alpha\beta L}{D} + \beta^2}.$$
(B.7)

B.7. Speed Reducer Design Optimization. The problem of designing a speed reducer [57] is a standard test problem. It consists of the design variables as face width x_1 , module of teeth x_2 , number of teeth on pinion x_3 , length of the first shaft between bearings x_4 , length of the second shaft between bearings x_5 , diameter of the first shaft x_6 , and diameter of the first shaft x_7 (all variables are continuous except x_3 that is integer). The weight of the speed reducer is to be minimized subject to constraints on bending stress of the gear teeth, surface stress, transverse deflections of the shafts, and stresses in the shaft [55]. The mathematical formulation of the problem, where $x^T = (x_1, x_2, x_3, x_4, x_5, x_6, x_7)$, is as follows:

Minimise
$$f(x) = 0.7854x_1x_2^2$$

$$\cdot \left(3.3333x_3^2 + 14.9334x_343.0934\right)$$

$$-1.508x_1\left(x_6^2 + x_7^3\right)$$

$$+7.4777\left(x_6^3 + x_7^3\right)$$

$$+0.7854\left(x_4x_6^2 + x_5x_7^2\right),$$
subject to $g_1(x) = \frac{27}{x_1x_2^2x_3} - 1 \le 0,$
 $g_2(x) = \frac{397.5}{x_1x_2^2x_3^2} - 1 \le 0,$
 $g_3(x) = \frac{1.93x_4^3}{x_2x_3x_6^4} - 1 \le 0,$
 $g_4(x) = \frac{1.93x_5^3}{x_2x_3x_7^4} - 1 \le 0,$
 $g_5(x) = \frac{1.0}{110x_6^3} \sqrt{\left(\frac{745.0x_4}{x_2x_3}\right)^2 + 16.9 \times 10^6}$

$$-1 \le 0,$$

$$g_{6}(x) = \frac{1.0}{85x_{7}^{3}} \sqrt{\left(\frac{745.0x_{5}}{x_{2}x_{3}}\right)^{2} + 157.5 \times 10^{6}}$$

$$-1 \le 0,$$

$$g_{7}(x) = \frac{x_{2}x_{3}}{40} - 1 \le 0,$$

$$g_{8}(x) = \frac{5x_{2}}{x_{1}} - 1 \le 0,$$

$$g_{9}(x) = \frac{x_{1}}{12x_{2}} - 1 \le 0,$$

$$g_{10}(x) = \frac{1.5x_{6} + 1.9}{x_{4}} - 1 \le 0,$$

$$g_{11}(x) = \frac{1.1x_{7} + 1.9}{x_{5}} - 1 \le 0.$$
(B.8)

The simple limits on the design variables are $2.6 \le x_1 \le 3.6$, $0.7 \le x_2 \le 0.8$, $17 \le x_3 \le 28$, $7.3 \le x_4 \le 8.3$, $7.8 \le x_5 \le 8.3$, $2.9 \le x_6 \le 3.9$, and $5.0 \le x_7 \le 5.5$.

B.8. Spring Design Optimisation. The main objective of this problem [58, 59] is to minimize the weight of a tension/compression string, subject to constraints of minimum deflection, shear stress, surge frequency, and limits on outside diameter and on design variables. There are three design variables: the wire diameter x_1 , the mean coil diameter x_2 , and the number of active coils x_3 [55]. The mathematical formulation of this problem, where $x^T = (x_1, x_2, x_3)$, is as follows:

Minimize
$$f(x) = (x_3 + 2) x_2 x_1^2$$
,
subject to $g_1(x) = 1 - \frac{x_2^3 x_3}{7,178 x_1^4} \le 0$,

$$g_2(x) = \frac{4x_2^2 - x_1 x_2}{12,566 (x_2 x_1^3) - x_1^4} + \frac{1}{5,108 x_1^2}$$

$$-1 \le 0$$
,

$$g_3(x) = 1 - \frac{140.45 x_1}{x_2^2 x_3} \le 0$$
,

$$g_4(x) = \frac{x_2 + x_1}{15} - 1 \le 0$$
.

The simple limits on the design variables are $0.05 \le x_1 \le 2.0$, $0.25 \le x_2 \le 1.3$, and $2.0 \le x_3 \le 15.0$.

Conflict of Interests

The authors declare that there is no conflict of interests regarding the publication of this paper.

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