

A New Construction of Zero Correlation Zone Sequences from Generalized Reed-Muller Codes

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Abstract—In this paper, we present a direct construction of zero-correlation zone (ZCZ) sequence sets (each associated with a graph) from the second order cosets of the first-order generalized Reed-Muller codes. This settles an open problem introduced by Rathinakumar and Chaturvedi in their 2008 paper.

I. INTRODUCTION

Zero-correlation zone (ZCZ) sequences were proposed in the late 1990s, aiming to eliminate the “near-far effect” and to achieve interference-free performance [1] in asynchronous code-division multiple access (CDMA) communication systems. By definition, ZCZ sequences possess zero auto- and cross- correlations over a region around the in-phase timing position. Hence, ZCZ sequence sets are used in a quasi-synchronous CDMA system, where received signals are approximately synchronized with transmitted signals [2].

The past decade has witnessed significant developments in the area of ZCZ sequence constructions [3]–[6]. In the literature, an interesting research direction has been to construct ZCZ sequences through a sequence transform of mutually orthogonal complementary codes (MOCCs). MOCCs can be represented as a set of two-dimensional matrices with zero aperiodic auto- and cross- correlation sums. For more details on perfect- and quasi- complementary codes, readers are referred to [7]–[16].

In [17], by constructing complete complementary codes (CCC) based on generalized Reed-Muller codes, Rathinakumar and Chaturvedi postulated that there may also exist direct constructions of ZCZ sequences using generalized Reed-Muller codes: “We believe that the direct construction of Golay complementary sets ... can be used to algebraically construct such ZCZ sets. One possible approach is to associate a graph with a ZCZ sequence set and relate the IFW¹ property of the ZCZ sequences to that graph.” This is referred to as the Rathinakumar-Chaturvedi open problem in this paper. Solving this open problem will uncover deeper inherent connections between ZCZ sequences, CCC, and Reed-Muller codes, which may be useful in finding new ZCZ sequence constructions. In addition, it will simplify the sequence generation and facilitate the sequence enumeration, owing to the Reed-Muller code properties.

¹i.e., interference free window.

Recently, Gong, Huo and Yang found that certain Golay sequences constructed from generalized Reed-Muller codes feature the zero autocorrelation zone (ZACZ) property [18]. In particular, they raised an open problem similar to that in [17]: how can CDMA spreading sequences be constructed from Golay sequences with the ZACZ property? A possible solution is to construct CDMA spreading sequences with zero auto- and cross- correlation zone properties from generalized Reed-Muller codes.

In this paper, we settle the Rathinakumar-Chaturvedi open problem by constructing ZCZ sequence sets from the second order cosets of the first-order generalized Reed-Muller codes. Interestingly, each of the proposed ZCZ sequence sets is associated with a two-layer graph where the ZCZ width is equal to the power of the vertex number in the upper layer of its graph.

II. PRELIMINARIES

Given two length- N complex-valued sequences $\mathbf{a} = [a_0, a_1, \dots, a_{N-1}]$ and $\mathbf{b} = [b_0, b_1, \dots, b_{N-1}]$, their aperiodic correlation function at time-shift τ is defined as

$$\rho(\mathbf{a}, \mathbf{b})(\tau) = \sum_{t=0}^{N-1-\tau} a_t b_{t+\tau}^*, \quad 0 \leq \tau \leq N-1. \quad (1)$$

When $\mathbf{a} \neq \mathbf{b}$, $\rho(\mathbf{a}, \mathbf{b})(\tau)$ is called the aperiodic cross-correlation function (ACCF); Otherwise, it is called aperiodic auto-correlation function (AACF). For simplicity, the AACF of $\mathbf{a}$ will be sometimes written as $\rho(\mathbf{a})(\tau)$.

In addition, the periodic cross- (or auto-) correlation function between $\mathbf{a}$ and $\mathbf{b}$ is defined as

$$\theta(\mathbf{a}, \mathbf{b})(\tau) = \rho(\mathbf{a}, \mathbf{b})(\tau) + \rho^*(\mathbf{b}, \mathbf{a})(N - \tau). \quad (2)$$

A. ZCZ sequences

Let $\mathcal{A} = \{\mathbf{a}_0, \mathbf{a}_1, \dots, \mathbf{a}_{K-1}\}$ be a set of K sequences, each of length L , i.e.,

$$\mathbf{a}_i = [a_0^i, a_1^i, \dots, a_n^i, \dots, a_{L-1}^i], \quad 0 \leq i \leq K-1. \quad (3)$$

Definition 1: $\mathcal{A}$ is said to be a (K, L, Z) -ZCZ sequence set if and only if it satisfies the following two conditions:

- 1) $\theta(\mathbf{a}_i)(\tau) = 0$ holds for any $0 \leq i \leq K-1$ and $1 \leq |\tau| \leq Z$;

2) $\theta(\mathbf{a}_i, \mathbf{a}_j)(\tau) = 0$ holds for any $i \neq j$ and $0 \leq |\tau| \leq Z$.

B. Complete Complementary Codes (CCC)

Let $\mathcal{C} = \{C_0, C_1, \dots, C_{K-1}\}$ be a set of K matrices, each of order $M \times N$, i.e.,

$$C_\mu = \begin{bmatrix} \mathbf{c}_0^\mu \\ \mathbf{c}_1^\mu \\ \vdots \\ \mathbf{c}_{M-1}^\mu \end{bmatrix}_{M \times N}, \quad 0 \leq \mu \leq K-1. \quad (4)$$

The aperiodic cross-correlation function between C_μ and C_ν (where $0 \leq \mu, \nu \leq K-1$), which is in the form of “aperiodic correlation sum”, is defined as

$$\rho(C_\mu, C_\nu)(\tau) = \sum_{m=0}^{M-1} \rho(\mathbf{c}_m^\mu, \mathbf{c}_m^\nu)(\tau). \quad (5)$$

Definition 2: $\mathcal{C}$ is said to be a (K, M, N) -mutually orthogonal complementary code set (MOCCS) if $\rho(C_\mu, C_\nu)(\tau) = 0$ for any $\mu \neq \nu$, and $\mu = \nu, \tau \neq 0$. It is known that $K \leq M$ [13]. When $K = M$, $\mathcal{C}$ is said to be a set of CCC.

C. Generalized Reed-Muller Codes for CCC Constructions

The following notations will be used throughout this paper.

- $\mathbb{Z}_q = \{0, 1, \dots, q-1\}$;
- $\xi = \exp(2\pi\sqrt{-1}/q)$, the q th root of unity;
- Denote by $[i_0, i_1, \dots, i_{n-1}]_2$ the binary representation of the integer $i = \sum_{\alpha=0}^{n-1} i_\alpha 2^\alpha$, where the most significant bit (MSB) is represented by i_{n-1} . For simplicity, we write $i = [i_0, i_1, \dots, i_{n-1}]_2$.
- Denote by $\bar{x} = 1 - x$ the binary complement of $x \in \{0, 1\}$.

Definition 3: For $\mathbf{x} = [x_0, x_1, \dots, x_{n-1}] \in \{0, 1\}^n$, the q -ary generalized Boolean function (GBF) $f(x_0, x_1, \dots, x_{n-1})$ is defined as a mapping $f : \{0, 1\}^n \rightarrow \mathbb{Z}_q$. The algebraic normal form of f can be written as a sum of monomials, i.e.,

$$f(x_0, x_1, \dots, x_{n-1}) = \sum_{0 \leq i \leq 2^n-1} c_i \prod_{\alpha=0}^{n-1} x_\alpha^{i_\alpha}, \quad c_i \in \mathbb{Z}_q. \quad (6)$$

If $r = \max\{\sum i_\alpha : c_i \sum i_\alpha \neq 0\}$, we say that the algebraic degree of f is r , i.e., $\deg(f) = r$. For $f_i = f(i_0, i_1, \dots, i_{n-1})$, define the $\mathbb{Z}_q$ -valued vector $\mathbf{f}$ and the complex-valued vector $\psi(f)$ (both of which are associated with f)

$$\begin{aligned} \mathbf{f} &= [f_0, f_1, \dots, f_{2^n-1}], \\ \psi(f) &= [\xi^{f_0}, \xi^{f_1}, \dots, \xi^{f_{2^n-1}}] \end{aligned} \quad (7)$$

respectively.

Definition 4: Denote by $\text{RM}_q(r, n)$ the r th order $\mathbb{Z}_q$ -valued generalized Reed-Muller code (which is a linear code over $\mathbb{Z}_q$) whose generator matrix is formed by the binary vectors associated with all monomials of algebraic degree at most r in $x_0, x_1, \dots, x_{n-1}$.

For instance, the following matrix

$$\begin{bmatrix} 11111111 \\ 01010101 \\ 00110011 \\ 00001111 \\ 00010001 \\ 00000101 \\ 00000011 \end{bmatrix} \begin{matrix} 1 \\ x_0 \\ x_1 \\ x_2 \\ x_0 x_1 \\ x_0 x_2 \\ x_1 x_2 \end{matrix}$$

is the generator matrix for $\text{RM}_4(2, 3)$ over $\mathbb{Z}_4$.

Definition 5: $\tilde{f} = f(\bar{x}_0, \bar{x}_1, \dots, \bar{x}_{n-1})$ is called the reversal of f because

$$\psi(\tilde{f}) = [\xi^{f_{2^n-1}}, \xi^{f_{2^n-2}}, \dots, \xi^{f_0}].$$

Definition 6: Let $J = \{j_0, j_1, \dots, j_{k-1}\} \subset \{0, 1, \dots, n-1\}$ and $\mathbf{x}_J = [x_{j_0}, x_{j_1}, \dots, x_{j_{k-1}}]$. $f|_{\mathbf{x}_J=\mathbf{d}}$ is called a restriction of f over constant $\mathbf{d} \in \{0, 1\}^k$. Specifically, $\psi(f|_{\mathbf{x}_J=\mathbf{d}})$ is a sequence of length 2^n with the elements coincide with that of $\psi(f)$ at the positions i where $i_{j_\alpha} = d_\alpha$ for each $0 \leq \alpha < k$, and equal to zero otherwise.

Next, we show existing constructions of complementary codes based on the second order cosets of $\text{RM}_q(1, m)$.

Lemma 1: (Construction of complementary codes [19])

Let $\deg(f) = 2$. Denote by $Q : \{0, 1\}^n \rightarrow \mathbb{Z}_q$ the quadratic part of f , where

$$Q = \sum_{0 \leq u < v \leq n-1} q_{uv} x_u x_v, \quad q_{uv} \in \mathbb{Z}_q. \quad (8)$$

Furthermore, denote by $G(Q)$ the graph of Q which is formed by joining vertices u and v by an edge if it has a non-zero weight of q_{uv} for any $0 \leq u < v \leq n-1$. Also, let $I = \{i_0, i_1, \dots, i_{n-k-1}\} \subset \{0, 1, \dots, n-1\} \setminus J$, where J is defined in Definition 6. Suppose that deleting the vertex set J from $G(Q)$ results in a Hamiltonian path, i.e.,

$$f|_{\mathbf{x}_J=\mathbf{d}} = \frac{q}{2} \sum_{\alpha=0}^{n-k-2} x_{i_{\pi(\alpha)}} x_{i_{\pi(\alpha+1)}} + \sum_{\alpha=0}^{n-k-1} l_\alpha x_{i_\alpha} + l, \quad (9)$$

where $l_\alpha, l \in \mathbb{Z}_q$, and π is a permutation of $(0, 1, \dots, n-k-1)$. Then

$$\left\{ f + \frac{q}{2} \left(\sum_{\alpha=0}^{k-1} y_\alpha x_{j_\alpha} + y_k x_a \right) : y_k, y_\alpha \in \mathbb{Z}_2 \right\} \quad (10)$$

generate a complementary matrix of order $2^{k+1} \times 2^n$ over $\mathbb{Z}_q$, where $a = i_{\pi(0)}$ or $i_{\pi(n-k-1)}$.

The following CCC construction is an extension of the construction in Lemma 1:

Lemma 2: (Construction of CCC [17])

Consider the q -ary GBF f in Lemma 1, and its reversal $\tilde{f}$. Let $t = [t_0, t_1, \dots, t_{k-1}]_2$. Define complementary codes C_t

and $C_{2^{k+t}}$ as

$$\left\{ f + \frac{q}{2} \left(\sum_{\alpha=0}^{k-1} y_{\alpha} x_{j_{\alpha}} + y_k x_a + \sum_{\gamma=0}^{k-1} t_{\gamma} x_{j_{\gamma}} \right) : y_k, y_{\alpha} \in \mathbb{Z}_2 \right\}, \quad (11)$$

and

$$\left\{ \tilde{f} + \frac{q}{2} \left(\sum_{\alpha=0}^{k-1} y_{\alpha} \bar{x}_{j_{\alpha}} + \bar{y}_k x_a + \sum_{\gamma=0}^{k-1} t_{\gamma} \bar{x}_{j_{\gamma}} \right) : y_k, y_{\alpha} \in \mathbb{Z}_2 \right\}, \quad (12)$$

respectively. Then

$$\{\psi(C_t) : 0 \leq t \leq 2^k - 1\} \cup \{\psi^*(C_{2^{k+t}}) : 0 \leq t \leq 2^k - 1\}, \quad (13)$$

generate a set of CCC, where $(\cdot)^*$ denotes the complex conjugate of $(\cdot)$.

III. PROPOSED ZCZ CONSTRUCTION FROM GENERALIZED REED-MULLER CODES

First, we need the following lemma.

Lemma 3: Let $q = 2$. For $k + 2$ binary variables in $x_0, x_1, \dots, x_k, x_{k+1}$, consider the following Boolean function

$$h = \sum_{\alpha=1}^{k+1} c_{\alpha} x_{\alpha} x_0 + \sum_{1 \leq \mu < \nu \leq k} d_{\mu\nu} x_{\mu} x_{\nu} + \sum_{\beta=0}^{k+1} e_{\beta} x_{\beta} + e', \quad (14)$$

where $c_{k+1} = 1, c_{\alpha} \in \mathbb{Z}_2$ for $1 \leq \alpha \leq k, d_{\mu\nu}, e_{\beta}, e' \in \mathbb{Z}_2$. Denote by $\mathbf{h}$ the resultant binary vector over $\mathbb{Z}_2^{k+2}$, i.e.,

$$\mathbf{h} = [h_0, h_1, \dots, h_{2^{k+2}-1}]. \quad (15)$$

For $0 \leq \tau \leq 2^{k+1} - 1$, we have

$$(-1)^{h_{\tau} + h_{\tau+1}} + (-1)^{h_{\tau+2^{k+1}} + h_{\tau+1+2^{k+1}}} = 0, \quad (16)$$

where the additions in the subscripts are performed modulo 2^{k+2} .

Proof: The proof of (16) follows if we can show

$$h_{\tau} + h_{\tau+1} + h_{\tau+2^{k+1}} + h_{\tau+1+2^{k+1}} \equiv 1 \pmod{2}, \quad (17)$$

for $0 \leq \tau \leq 2^{k+1} - 1$. For $0 \leq \tau \leq 2^{k+1} - 2$, denote by

$$\begin{aligned} \tau &= [\tau_0, \tau_1, \dots, \tau_k, 0]_2, \\ \tau + 1 &= [\tau'_0, \tau'_1, \dots, \tau'_k, 0]_2. \end{aligned} \quad (18)$$

Clearly,

$$\begin{aligned} \tau + 2^{k+1} &= [\tau_0, \tau_1, \dots, \tau_k, 1]_2, \\ \tau + 1 + 2^{k+1} &= [\tau'_0, \tau'_1, \dots, \tau'_k, 1]_2. \end{aligned} \quad (19)$$

By (14), (18) and (19), we have

$$\begin{cases} h_{\tau} &= \sum_{\alpha=1}^{k+1} c_{\alpha} \tau_{\alpha} \tau_0 + \sum_{1 \leq \mu < \nu \leq k} d_{\mu\nu} \tau_{\mu} \tau_{\nu} \\ &+ \sum_{\beta=0}^{k+1} e_{\beta} \tau_{\beta} + e', \\ h_{\tau+2^{k+1}} &= h_{\tau} + \tau_0 + e_{k+1} \end{cases} \quad (20)$$

and

$$\begin{cases} h_{\tau'} &= \sum_{\alpha=1}^{k+1} c_{\alpha} \tau'_{\alpha} \tau'_0 + \sum_{1 \leq \mu < \nu \leq k} d_{\mu\nu} \tau'_{\mu} \tau'_{\nu} \\ &+ \sum_{\beta=0}^{k+1} e_{\beta} \tau'_{\beta} + e', \\ h_{\tau'+2^{k+1}} &= h_{\tau'} + \tau'_0 + e_{k+1} \end{cases} \quad (21)$$

Therefore,

$$\begin{aligned} &h_{\tau} + h_{\tau+1} + h_{\tau+2^{k+1}} + h_{\tau+1+2^{k+1}} \\ &= \underbrace{h_{\tau} + h_{\tau+2^{k+1}}}_{\equiv \tau_0 + \tau'_0} + \underbrace{h_{\tau+1} + h_{\tau+1+2^{k+1}}}_{\equiv 1 \pmod{2}} \\ &\equiv 1 \pmod{2}, \end{aligned}$$

which is (17), thus completing the proof. $\blacksquare$

Example 1: Let $k = 2$. The following Boolean function

$$h = x_3 x_0 + x_0 x_2 + x_2 x_1 + x_3 + x_0 + 1 \quad (22)$$

results in

$$\mathbf{h} = [\underbrace{1, 0, 1, 0, 1, 1, 0, 0}_{\text{length}=8}, \underbrace{0, 0, 0, 0, 0, 1, 1, 0}_{\text{length}=8}]$$

which satisfies (16). In particular, if $h = x_3 x_0$, one gets

$$\mathbf{h} = [\underbrace{0, 0, 0, 0, 0, 0, 0, 0}_{\text{length}=8}, \underbrace{0, 1, 0, 1, 0, 1, 0, 1}_{\text{length}=8}]$$

which also satisfies (16).

Theorem 1: Consider a set of CCC generated by *Lemma 2*, i.e., $\mathcal{C} = \{C_0, C_1, \dots, C_{2^{k+1}-1}\}$, where

$$C_{\mu} = \begin{bmatrix} \mathbf{c}_0^{\mu} \\ \mathbf{c}_1^{\mu} \\ \vdots \\ \mathbf{c}_{2^{k+1}-1}^{\mu} \end{bmatrix}_{2^{k+1} \times 2^n}, \quad 0 \leq \mu \leq 2^{k+1} - 1. \quad (23)$$

Based on $\mathbf{h}$ in (15), define two matrices $\mathcal{H}^0$ and $\mathcal{H}^1$ in (24) and (25), respectively. Then, the rows of matrix

$$\mathcal{H} = [\mathcal{H}^0, \mathcal{H}^1]$$

form a

$$(K = 2^{k+1}, L = 2^{n+k+2}, Z = 2^n)$$

ZCZ sequence set.

Proof: Denote by $\mathcal{H}_i$ the i th row vector of $\mathcal{H}$. For $0 \leq \tau \leq 2^n$, the periodic cross-correlation function between $\mathcal{H}_i$ and $\mathcal{H}_j$ is shown in (26).

Next, by (16), (26) and the zero aperiodic sum property of $\mathcal{C}$, we have

$$\begin{aligned} &\theta(\mathcal{H}_i, \mathcal{H}_j)(\tau) \\ &= 2 \cdot \sum_{m=0}^{2^{k+1}-1} \rho(\mathbf{c}_m^i, \mathbf{c}_m^j)(\tau) \\ &= 2\rho(C_i, C_j)(\tau) \\ &= \begin{cases} 2^{k+n+2}, & \text{if } i = j \text{ and } \tau = 0; \\ 0, & \text{otherwise} \end{cases} \end{aligned} \quad (27)$$

$$\mathcal{H}^0 = \begin{bmatrix} \mathbf{c}_0^0(-1)^{h_0}, & \mathbf{c}_1^0(-1)^{h_1}, & \cdots, & \mathbf{c}_{2^{k+1}}^0(-1)^{h_{2^{k+1}-1}} \\ \mathbf{c}_0^1(-1)^{h_0}, & \mathbf{c}_1^1(-1)^{h_1}, & \cdots, & \mathbf{c}_{2^{k+1}}^1(-1)^{h_{2^{k+1}-1}} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{c}_0^{2^{k+1}-1}(-1)^{h_0}, & \mathbf{c}_1^{2^{k+1}-1}(-1)^{h_1}, & \cdots, & \mathbf{c}_{2^{k+1}}^{2^{k+1}-1}(-1)^{h_{2^{k+1}-1}} \end{bmatrix}. \quad (24)$$

$$\mathcal{H}^1 = \begin{bmatrix} \mathbf{c}_0^0(-1)^{h_{2^{k+1}}}, & \mathbf{c}_1^0(-1)^{h_{2^{k+1}+1}}, & \cdots, & \mathbf{c}_{2^{k+1}}^0(-1)^{h_{2^{k+2}-1}} \\ \mathbf{c}_0^1(-1)^{h_{2^{k+1}}}, & \mathbf{c}_1^1(-1)^{h_{2^{k+1}+1}}, & \cdots, & \mathbf{c}_{2^{k+1}}^1(-1)^{h_{2^{k+2}-1}} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{c}_0^{2^{k+1}-1}(-1)^{h_{2^{k+1}}}, & \mathbf{c}_1^{2^{k+1}-1}(-1)^{h_{2^{k+1}+1}}, & \cdots, & \mathbf{c}_{2^{k+1}}^{2^{k+1}-1}(-1)^{h_{2^{k+2}-1}} \end{bmatrix}. \quad (25)$$

$$\begin{aligned} & \theta(\mathcal{H}_i, \mathcal{H}_j)(\tau) \\ &= 2 \cdot \sum_{m=0}^{2^{k+1}-1} \rho(\mathbf{c}_m^i, \mathbf{c}_m^j)(\tau) + \left[\underbrace{(-1)^{h_{2^{k+1}-1}+h_{2^{k+1}}} + (-1)^{h_{2^{k+2}-1}+h_0}}_{\text{}} \right] \rho^*(\mathbf{c}_0^j, \mathbf{c}_{2^{k+1}-1}^i)(L-\tau) \\ &+ \sum_{m=0}^{2^{k+1}-2} \left[\underbrace{(-1)^{h_m+h_{m+1}} + (-1)^{h_{m+2^{k+1}}+h_{m+1+2^{k+1}}}}_{\text{}} \right] \rho^*(\mathbf{c}_{m+1}^j, \mathbf{c}_m^i)(L-\tau). \end{aligned} \quad (26)$$

Thus, the proof of *Theorem 1* follows. $\blacksquare$

Remark 1: The proposed ZCZ construction in *Theorem 1* is a generalization of that in [3], where

$$\mathbf{h} = [\underbrace{0, 0, \dots, 0, 0}_{\text{length}=2^{k+1}}, \underbrace{0, 1, \dots, 0, 1}_{\text{length}=2^{k+1}}]$$

(which is generated by $h = x_0 x_{k+1}$) is used, as indicated by the latter case of *Example 1*.

Based on *Lemma 2*, *Lemma 3* and *Theorem 1*, we are now ready to present the proposed ZCZ construction based on generalized Reed-Muller codes. For completeness, some definitions will be given again.

Theorem 2: For $n+k+2$ binary variables $x_0, x_1, \dots, x_{n+k+1}$, define a GBF $f(x_0, x_1, \dots, x_{n-1})$ which satisfies the condition in *Lemma 1*, i.e., deleting the vertex set $J = \{j_0, j_1, \dots, j_{k-1}\}$ from the graph of f results in a Hamiltonian path characterized by $\frac{q}{2} \sum_{\alpha=0}^{n-k-2} x_{i_{\pi(\alpha)}} x_{i_{\pi(\alpha+1)}}$ over $I = \{i_0, i_1, \dots, i_{n-k-1}\}$, where π is a permutation of $(0, 1, \dots, n-k-1)$, $I \cup J = \{0, 1, \dots, n-1\}$ and $I \cap J = \emptyset$. Also, define another GBF $h(x_n, x_{n+1}, \dots, x_{n+k+1})$ as

$$\begin{aligned} h &= \frac{q}{2} \sum_{\alpha=1}^{k+1} c_\alpha x_{\alpha+n} x_n + \frac{q}{2} \sum_{1 \leq \mu < \nu \leq k} d_{\mu\nu} x_{\mu+n} x_{\nu+n} \\ &+ \frac{q}{2} \sum_{\beta=0}^{k+1} e_\beta x_{\beta+n}, \end{aligned} \quad (28)$$

where $c_{k+1} = 1, c_\alpha \in \mathbb{Z}_2$ for $1 \leq \alpha \leq k$, $d_{\mu\nu}, e_\beta \in \mathbb{Z}_2$. Let ϖ be a permutation of $(0, 1, \dots, k)$, and $t = [t_0, t_1, \dots, t_{k-1}]_2$. Also, let a be either end-vertex of the Hamiltonian path, i.e., $a = i_{\pi(0)}$ or $i_{\pi(n-k-1)}$. Generate two sequence sets $\{\mathbf{z}_t : 0 \leq$

$t \leq 2^k - 1\}$ and $\{\mathbf{z}_{t+2^k} : 0 \leq t \leq 2^k - 1\}$ by

$$\left\{ f + h + \frac{q}{2} \left(\sum_{\alpha=0}^{k-1} x_{n+\varpi(\alpha)} x_{j_\alpha} + x_{n+\varpi(k)} x_a + \sum_{\gamma=0}^{k-1} t_\gamma x_{j_\gamma} \right) \right\} \quad (29)$$

and

$$\left\{ \tilde{f} + h + \frac{q}{2} \left(\sum_{\alpha=0}^{k-1} x_{n+\varpi(\alpha)} \bar{x}_{j_\alpha} + \bar{x}_{n+\varpi(k)} x_a + \sum_{\gamma=0}^{k-1} t_\gamma \bar{x}_{j_\gamma} \right) \right\} \quad (30)$$

respectively. Then,

$$\{\psi(\mathbf{z}_t)\} \cup \{\psi^*(\mathbf{z}_{t+2^k})\}, \quad \text{where } 0 \leq t \leq 2^k - 1 \quad (31)$$

form a $(K = 2^{k+1}, L = 2^{n+k+2}, Z = 2^n)$ -ZCZ sequence set.

Proof: The proof of this theorem can be obtained by applying *Lemma 2* and *Theorem 1*. It is however omitted due to the space limitation. $\blacksquare$

Example 2: Let $q = 4, k = 1, n = 3, J = \{0\}$ and $I = \{1, 2\}$. Also, let

$$\begin{aligned} f &= 2x_1 x_2 + 3x_0 x_2 + x_0 x_1 + 3x_1 + x_2 + 1, \\ h &= 2x_3 x_5 + 2x_3 x_4 + 2x_3 + 2x_5 + 3. \end{aligned}$$

Clearly,

$$\begin{aligned} \tilde{f} &= f(1 - x_0, 1 - x_1, 1 - x_2) \\ &= 2x_1 x_2 + 3x_0 x_2 + x_0 x_1 + 2x_1 + 2x_2 + 3. \end{aligned}$$

Generate two sequence sets $\{\mathbf{z}_t : t \in \mathbb{Z}_2\}$ and $\{\mathbf{z}_{t+2} : t \in \mathbb{Z}_2\}$ by

$$\left\{ f + h + 2(x_0 x_3 + x_1 x_4) + 2t x_0 \right\}$$

and

$$\left\{ \tilde{f} + h + 2(\bar{x}_0 x_6 + x_1 \bar{x}_4) + 2t \bar{x}_0 \right\}$$

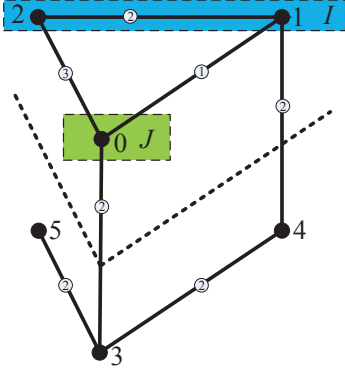


Fig. 1. Graph representation of $f + h + 2(x_0x_3 + x_1x_4)$

respectively. A graph representation of $f + h + 2(x_0x_3 + x_1x_4)$ is given in Fig. 1, with every edge's weight is shown in a small circle. It is seen that this graph features a two-layer structure where the upper layer and lower layer (separated by the dashed lines) are specified by vertex sets $\{0, 1, 2\}$ and $\{3, 4, 5\}$, respectively. Also, deleting the vertex set $J = \{0\}$, its upper layer graph reduces to a Hamiltonian path (characterized by $2x_1x_2$) over $I = \{1, 2\}$. For $\{z_0, z_1, z_2, z_3\}$ shown in (32), we by *Theorem 2* assert that

$$\{\xi^{z_0}, \xi^{z_1}\} \cup \{\xi^{-z_2}, \xi^{-z_3}\}$$

form a $(K = 4, L = 64, Z = 8)$ -ZCZ sequence set, where $\xi = \sqrt{-1}$. Interestingly, the ZCZ width is equal to the power of vertex number in the upper layer of the graph in Fig. 1.

$$\begin{aligned} z_0 &= [00301022201030020012100002101202 \\ &\quad 22123200201030022230322202101202], \\ z_1 &= [02321220221232000210120200121000 \\ &\quad 20103002221232002032302000121000], \\ z_2 &= [22230322202101202201030002212320 \\ &\quad 00012100202101200023212202212320], \\ z_3 &= [02032302000121000221232022010300 \\ &\quad 20210120000121002003010222010300]. \end{aligned} \quad (32)$$

IV. CONCLUSIONS

In this paper, we have settled the Rathinakumar-Chaturvedi open problem positively by presenting a new construction of ZCZ sequences from generalized Reed-Muller codes. As shown in *Example 2*, every proposed ZCZ sequence set is associated with a two-layer graph (see Fig. 1), where the ZCZ width is equal to the power of vertex number (i.e., 2^n) in the upper layer graph. Our proposed construction also extends and generalizes the ZCZ construction in [3], thus yielding many new ZCZ sequences which have not been discovered before, as shown in *Remark 1*.

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